

CS 370: INTRODUCTION TO SECURITY
04.06: CRYPTOGRAPHY BASICS

Tu/Th 4:00 – 5:50 PM (WNGR 149)

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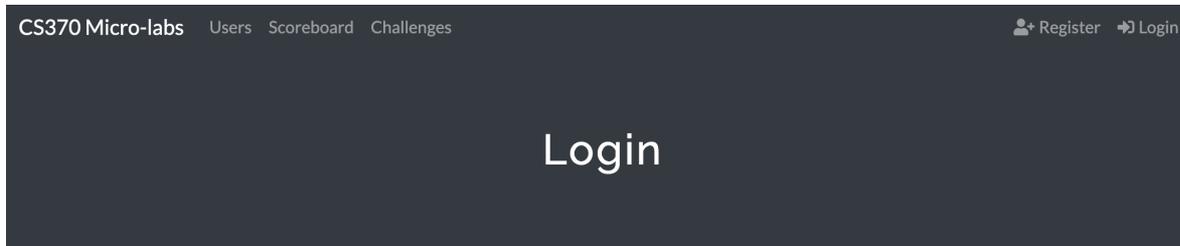
SAIL
Secure AI Systems Lab

TOPICS FOR TODAY

- Micro-lab instruction
- Crypto basics
 - Why do we need crypto?
 - What were the ancient crypto schemes?
 - What does it mean by **perfectly secure**?
 - What were the perfect crypto schemes so far?
 - What are the limitations of those above?
 - What were some practical solutions and their (also) limitations?

MICRO-LAB INSTRUCTION

- Create an account on ctf.secure-ai.systems
 - Use OSU email address
 - Use some secure password
 - Once registered, check the inbox for the welcome email
 - Otherwise, lmk



User Name or Email

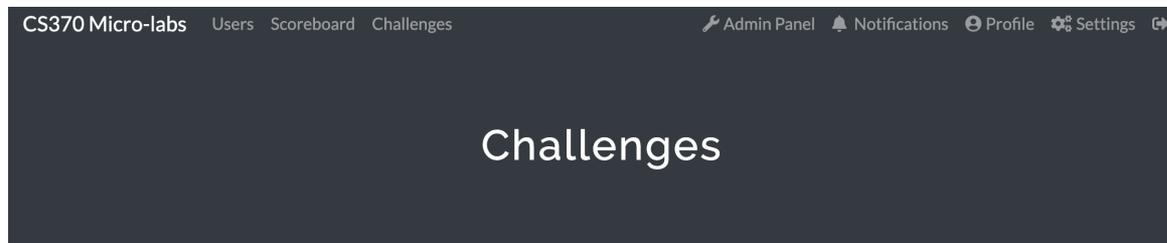
Password

[Forgot your password?](#)

[Submit](#)

MICRO-LAB INSTRUCTION – CONT'D

- Go to **Challenges**
 - You can find two challenges
 - More will come soon



Week0



MICRO-LAB INSTRUCTION – CONT'D

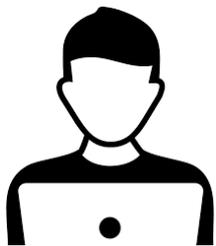
- Connect to the “Solve” server
 - This is the place where you can solve the challenges
 - Instruction can be found on here ([the course website](#))
 - Your username is set to your ONID
 - Password for logging in can be found in the Canvas announcement
 - I will walk you through how to do it now...

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WHY DO WE NEED CRYPTO?

- Confidentiality
 - We want to communicate with others securely (and privately)



Let's have Local Boyz for dinner!



WHY DO WE NEED CRYPTO?

- Confidentiality
 - We want to communicate with others securely (and privately)
 - Plaintext communication can be eavesdropped by an adversary



Let's have Local Boyz for dinner!

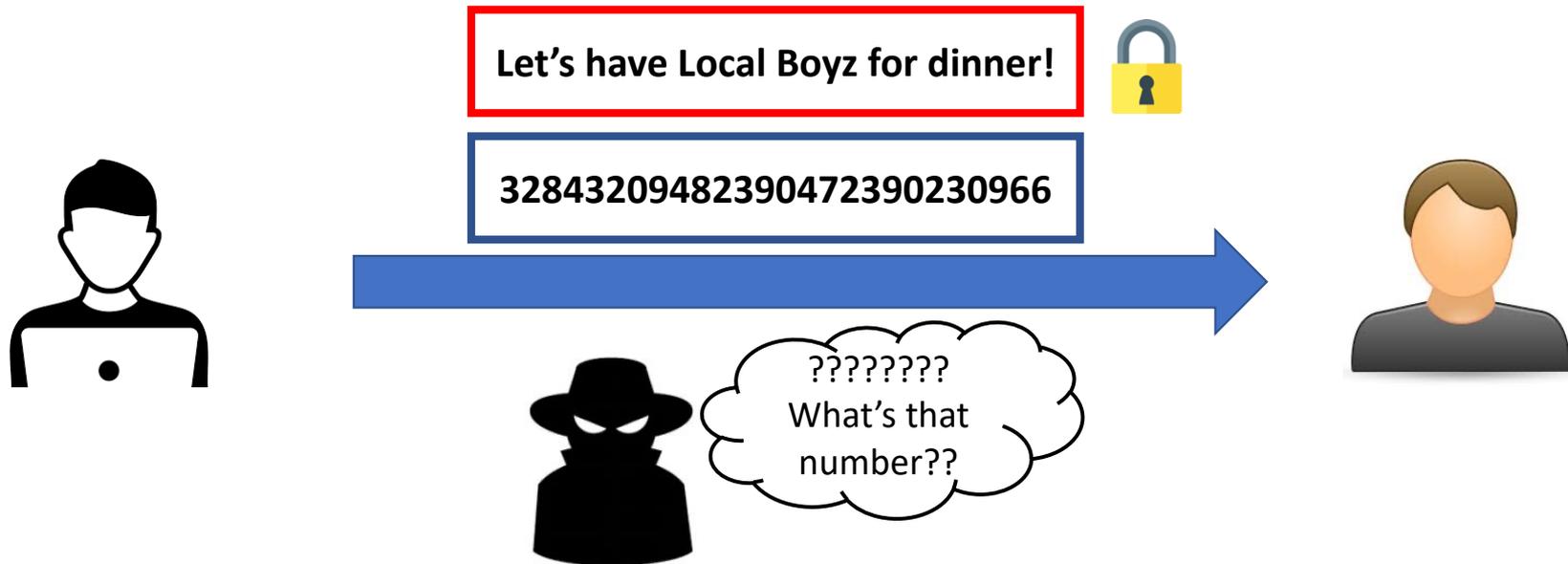


Local Boyz...

WHY DO WE NEED CRYPTO?

- Confidentiality

- We want to communicate with others securely (and privately)
- Plaintext communication can be eavesdropped by an adversary
- Cryptography enables secure (and private) communication



BASIC TERMINOLOGY

- Terms

- Plaintext: readable text, before getting encrypted
- Ciphertext: encrypted text, transformed plaintext using an encryption algorithm
- Encryption/decryption: the act of encrypting (or decrypting)



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CAESAR CIPHER

- Crypto scheme in Roman empire
 - Encryption: shift each character by N
 - Example: shift by 3
 - Plaintext: ABCDEFGHIJKLMNOPQRSTUVWXYZ
 - Cipher text: **DEFGHIJKLMNOPQRSTUVWXYZABC**

 - Plaintext: HELLO
 - Cipher text: **KHOOR**



ARE WE SAFE NOW?

PROBLEM(S) IN CAESAR CIPHER

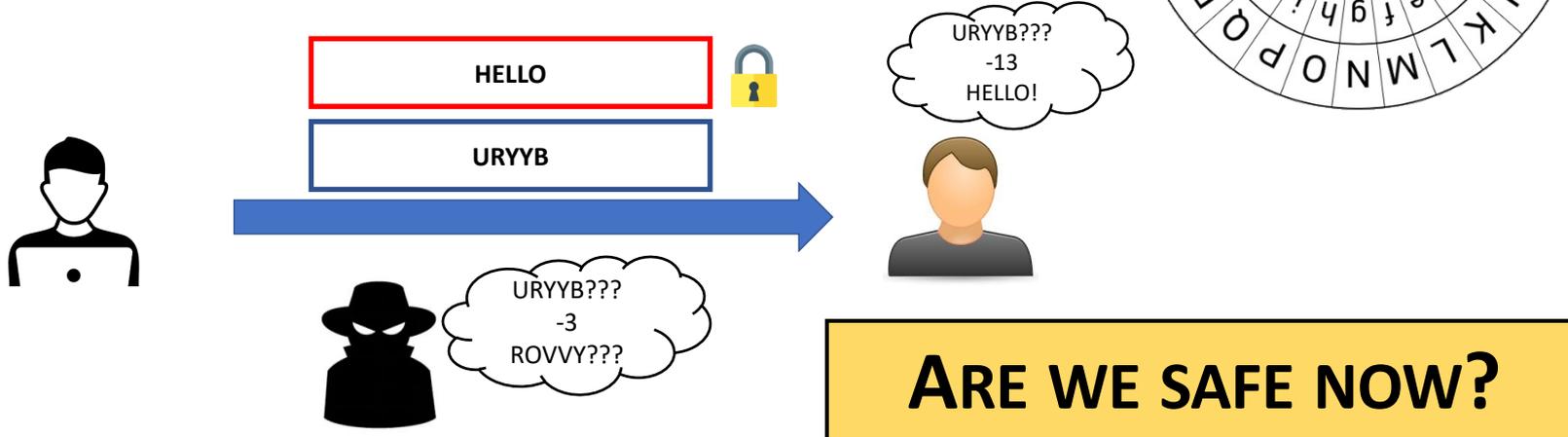
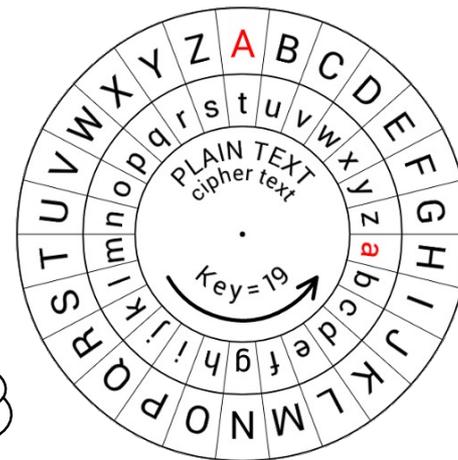
- What if:
 - An adversary knows the shift offset?



GENERAL FORM OF CAESAR CIPHER

- Rot-N cipher

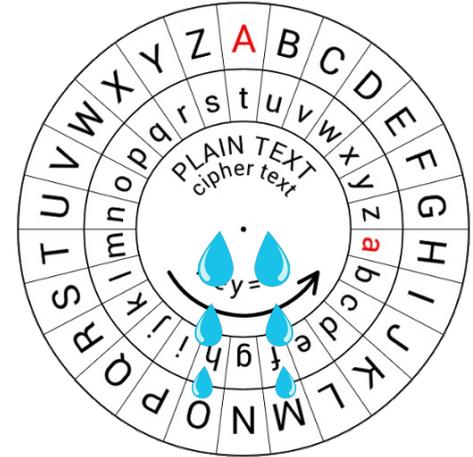
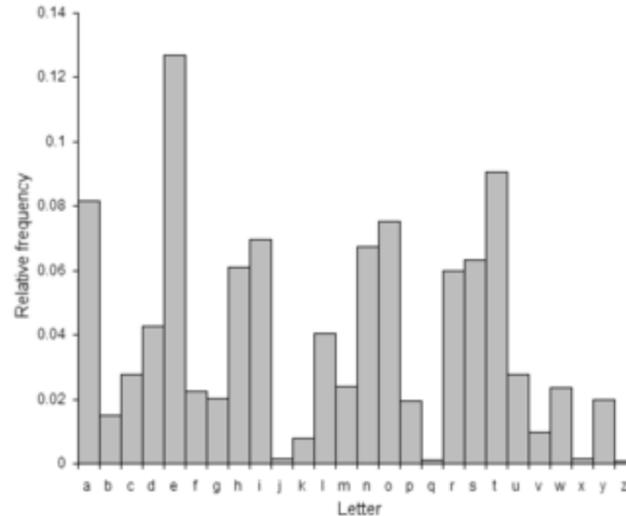
- Encryption: shift each character by N
- More complex than Caesar cipher
 - We can set N and share between us
 - If an adversary doesn't know N, it's hard to decrypt it
- Example: shift by 13



PROBLEM(S) IN ROT-N CIPHER

- What if:
 - An adversary knows the shift offset?
 - An attacker finds the offset

E	11.1607%	56.88	M	3.0129%	15.36
A	8.4966%	43.31	H	3.0034%	15.31
R	7.5809%	38.64	G	2.4705%	12.59
I	7.5448%	38.45	B	2.0720%	10.56
O	7.1635%	36.51	F	1.8121%	9.24
T	6.9509%	35.43	Y	1.7779%	9.06
N	6.6544%	33.92	W	1.2899%	6.57
S	5.7351%	29.23	K	1.1016%	5.61
L	5.4893%	27.98	V	1.0074%	5.13
C	4.5388%	23.13	X	0.2902%	1.48
U	3.6308%	18.51	Z	0.2722%	1.39
D	3.3844%	17.25	J	0.1965%	1.00
P	3.1671%	16.14	Q	0.1962%	(1)



Letter frequency in English

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PERFECT SECURITY

- Shannon's intuition
 - An adversary should not distinguish a message M from a random text R



Claude Shannon (1916 ~ 2001)
A Father of Information Theory
and Modern Cryptography

PERFECT SECURITY

- Shannon's intuition

- An adversary should not distinguish a message M from a random text R

- Formally:

- $\Pr[M = m | C = c] = \Pr[M = m]$

- where

- m is a message (from a set M)

- c is a ciphertext (from a set of all ciphertexts C)

- $\Pr[C = c | M = m] = \Pr[C = c]$

- It means:

- Ciphertext provides no additional information

- Observing c does not help with guessing $M = m$

- c is independent of the message m



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XOR CIPHER

- Crypto scheme with perfect secrecy
 - Encryption:
 - Given a message m and a random key k
 - Ciphertext $c = m \oplus k$
 - Example:
 - Message: HELLO
 - Key : ABCDE

A	B	Q
0	0	0
0	1	1
1	0	1
1	1	0

Message	H (0x48)	E (0x45)	L (0x4c)	L (0x4c)	O (0x4f)
Key	A (0x41)	B (0x42)	C (0x43)	D (0x44)	E (0x45)
Ciphertext	0x9	0x7	0xf	0x8	0xa

XOR CIPHER

- Crypto scheme with perfect secrecy
 - Encryption:
 - Given a message m and a random key k
 - Plaintext $m = k \oplus c$
 - Example:
 - Message: HELLO
 - Key : ABCDE

A	B	Q
0	0	0
0	1	1
1	0	1
1	1	0

Key	A (0x41)	B (0x42)	C (0x43)	D (0x44)	E (0x45)
Ciphertext	0x9	0x7	0xf	0x8	0xa
Decrypt	H	E	L	L	O

XOR CIPHER: IN BITWISE OPERATION

- Example from Wikipedia¹

The string "Wiki" (01010111 01101001 01101011 01101001 in 8-bit ASCII) can be encrypted with the repeating key 11110011 as follows:

$$\begin{array}{r} 01010111 \ 01101001 \ 01101011 \ 01101001 \\ \oplus 11110011 \ 11110011 \ 11110011 \ 11110011 \\ \hline = 10100100 \ 10011010 \ 10011000 \ 10011010 \end{array}$$

And conversely, for decryption:

$$\begin{array}{r} 10100100 \ 10011010 \ 10011000 \ 10011010 \\ \oplus 11110011 \ 11110011 \ 11110011 \ 11110011 \\ \hline = 01010111 \ 01101001 \ 01101011 \ 01101001 \end{array}$$

A	B	Q
0	0	0
0	1	1
1	0	1
1	1	0

¹Image from: https://en.wikipedia.org/wiki/XOR_cipher

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PROBLEM(S) IN XOR CIPHER

- What if:
 - An adversary accidentally knows a pair of m and c
 - Key $k = m \oplus c$

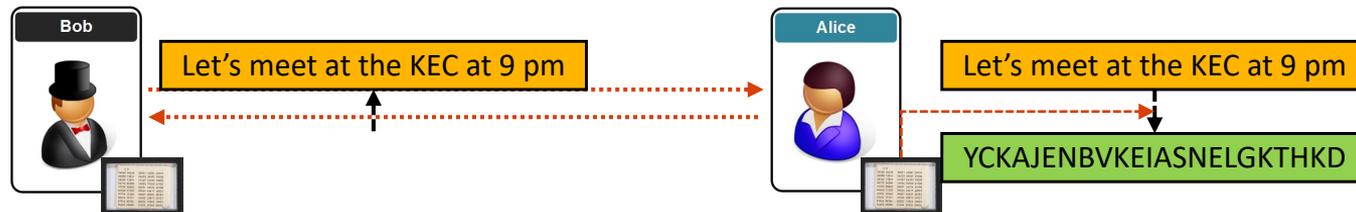
Message	H (0x48)	E (0x45)	L (0x4c)	L (0x4c)	O (0x4f)
Ciphertext	0x9	0x7	0xf	0x8	0xa
Key	A (0x41)	B (0x42)	C (0x43)	D (0x44)	E (0x45)

GENERAL FORM OF XOR CIPHER

- OTP
 - One-Time Pads (OTP) is an encryption mechanism
- How it works?
 - Alice and Bob want to communicate *securely*
 - Alice and Bob share the same OTP
 - Alice encrypts a message to send with the OTP
 - Alice sends the encrypted message to Bob
 - Bob decrypts the received message with the OTP



An Example OTP



PROBLEM(S) IN XOR CIPHER

- What if:
 - An adversary accidentally knows a pair of m and c
 - Key $k = m \oplus c$

Message	H (0x48)	E (0x45)	L (0x4c)	L (0x4c)	O (0x4f)
Ciphertext	0x9	0x7	0xf	0x8	0xa
Key	A (0x41)	B (0x42)	C (0x43)	D (0x44)	E (0x45)

- Practical limitations:
 - What if we want to encrypt a 1GB video file?
 - How can we share keys with others (OTP)?

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STREAM CIPHER

- Reduce key generation and exchange overheads
 - Encryption:
 - Given a message m and a random key k
 - Ciphertext $c = m \oplus k$
 - and:
 - The key stream is **generated by the same mechanism** for a sender and a receiver
 - The key stream is a **byte stream** (0xAB129dB...)
 - It **performs XOR encryption** over this byte stream

STREAM CIPHER

- Stream cipher
 - Example:

Encrypt message 1 with 0x1b395a46
Encrypt message 2 with 0xf1737202
Encrypt message 3 with 0xccf0de05...



A random number generator ...

1: 0x1b395a46
2: 0xf1737202
3: 0xccf0de05
4: 0x908b0feb
5: 0x9d4c9add

Decrypt message 1 with 0x1b395a46
Decrypt message 2 with 0xf1737202
Decrypt message 3 with 0xccf0de05...



A random number generator

1: 0x1b395a46
2: 0xf1737202
3: 0xccf0de05
4: 0x908b0feb
5: 0x9d4c9add
...

STREAM CIPHER

- Stream cipher
 - Example: [RC4](#)/[RC5](#)

Encrypt message 1 with 0x1b395a46
Encrypt message 2 with 0xf1737202
Encrypt message 3 with 0xccf0de05...

Decrypt message 1 with 0x1b395a46
Decrypt message 2 with 0xf1737202
Decrypt message 3 with 0xccf0de05...



Physically meet and share
40~2048 bits of keys
(max 512 bytes, short!)



RC4 or
RC5
Algorithm

1: 0x1b395a46
2: 0xf1737202
3: 0xccf0de05
4: 0x908b0feb
5: 0x9d4c9add
...

Key: 0xd7fe6798a7c6a9859b289ce

1: 0x1b395a46
2: 0xf1737202
3: 0xccf0de05
4: 0x908b0feb
5: 0x9d4c9add
...



RC4 or
RC5
Algorithm

Key: 0xd7fe6798a7c6a9859b289ce

PROBLEM(S) IN RC4/RC5

- See the Wikipedia sections
 - Bit-flipping attacks
 - Reused key attacks
 - Differential attacks
 - ...

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